**Chunk List**

Daniel J. Szelogowski

University of Wisconsin – Whitewater

© 2017

**PREFACE**

Chunking data is obviously no new concept; however, I had never found any data structures that used chunking as the basis of their implementation. I figured that by using chunking alongside concurrency, I could create an extremely fast run-time in regards to particular methods as searching and/or sorting.

By using chunking and concurrency to my advantage, I came up with the chunk list — a dynamic list-based data structure that would separate large amounts of data into specifically sized chunks, each of which should be able to be searched at the exact same time by searching each chunk on a separate thread.

As a result of implementing this concept into its own class, I was able to create something that almost consistently gives around 3.5x faster results than a regular ArrayList. However, should speed be a particular issue even after implementation, users can modify the size of the chunks and benchmark the speed of using smaller or larger chunks, depending on the amount of data being stored. Generally a chunk size of 1/10th the size of the data being stored gives the fastest and most consistent results.

Notes:

* A full implementation can be found at <https://github.com/danielathome19/Chunk-List>
* All code examples given are in the C# language.
* Examples are given for each main method the class should have implemented.
* Time complexities are listed for each method.

# TABLE OF CONTENTS

1. Discussion1

1.1 What is a Chunk List?1

1.1.1 Defining a Chunk1

1.1.2 Efficiency of Chunking1

1.2 Where is a Chunk List Used?2

1.3 Benefits2

1.3.1 Area-Specific Usage1

1.3.2 Competitive Usage1

2. Implementation Details3

2.1 Construction3

2.2 Multithreading Methods4

2.3 Index-Based Methods5

2.3.1 Index Issues6

2.3.2 Chunk Resizing9

2.4 Element Operations5

2.4.1 Adding Elements6

2.4.2 Removing Elements7

2.4.3 Searching8

2.4.4 Sorting11

3. Complexity Analysis12

3.1 Element-Based Methods12

3.2 List-Based Methods12

4. Usage & Examples12

4.1 Modern Usage12

4.1.1 Experiments6

4.1.2 Unit Test & Benchmarks6

4.2 Related Work12

References12

Appendix – Method Headers & Complexity Table12

**1. DISCUSSION**

**1.1 What is a Chunk List?**

A chunk list is an array-based list of elements in which data is stored in inner lists of a certain capacity, allowing for easily modifiable and faster runtimes based on the number of elements being stored. A simple way to conceptualize a chunk list would be an ArrayList (dynamic array) of ArrayLists.

The main list would contain the “chunks”, or ArrayLists that are not allowed to be filled past a specific capacity.

Any time a “chunk” has reached capacity, a new ArrayList is added and items are added to that chunk from thereon.

By doing this process and splitting our list into chunks, we can use parallel processing to our advantage. Using concurrency, we can run each chunk on a separate thread when doing tasks such as searching or removing.

This can be expressed visually as a table, in the case of a chunk list containing the numbers 1 – 50 where the chunk size is set to 10 elements. (See Figure 1)

***Fig. 1:***

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| 1 2 3 4 5 6 7 8 9 10 | 11 12 13 14 15 16 17 18 19 20 | 21 22 23 24 25 26 27 28 29 30 | 31 32 33 34 35 36 37 38 39 40 | 41 42 43 44 45 46 47 48 49 50 |

**1.1.1 Defining a Chunk**

A CHUNK LIST IS AN ARRAY-BASED LIST OF ELEMENTS IN WHICH

**1.1.2 Efficiency of Chunking**

A CHUNK LIST IS AN ARRAY-BASED LIST OF ELEMENTS IN WHICH

**1.2 Where is a Chunk List Used?**

Useful for storing very large and very small amounts of elements.

Benefits shine especially when list is unsorted:

* Fast searching
* Fast removal
* Fast insertion

Implementation is easy and short.

Sorting is quick even with large amounts of chunks.

Timed testing results typically yield to be approximately 3.5x faster than an ArrayList on average.

In any scenario, a chunk list can be used in place of an ArrayList especially, as well as something such as a Binary Search Tree, as searching may be faster based on processing power.

**1.3 Benefits**

A CHUNK LIST IS AN ARRAY-BASED LIST OF ELEMENTS IN WHICH

**1.3.1 Area-Specific Usage**

A CHUNK LIST IS AN ARRAY-BASED LIST OF ELEMENTS IN WHICH

**1.3.2 Competitive Usage**

A CHUNK LIST IS AN ARRAY-BASED LIST OF ELEMENTS IN WHICH

**2. IMPLEMENTATION**

**2.1 Construction**

The basis of the chunk list is the inner list. This is best implemented using some sort of dynamic list, such as ArrayList (or List in C#).

This inner list will start out with a single list on the inside.

Constructor must include an integer, chunk size. Otherwise, revert to a default size.

New lists (chunks) will only be added to the main list when the chunk at the end has reached capacity.

A chunk list may be implemented with generics (or templates) so long as the generic type is comparable.

*Example:*

using System;

using Sytem.Collections.Generic;

using System.Threading.Tasks;

class ChunkList<T> where T : IComparable

{

private List<List<T>> myList;

private int chunkSize;

private const int DEFAULT\_SIZE = 1000;

public ChunkList() : this(DEFAULT\_SIZE)

{

}

public ChunkList(int chunkSize)

{

this.chunkSize = chunkSize;

myList = new List<List<T>>();

}

}

## 2.2 Multithreading Methods

Multithreading is an especially important part of chunk list implementation, as the basis of the list’s speed is primarily the result of concurrency. For most methods in a chunk list, a new thread can be created for each chunk to be iterated through.

A good example of this lies within C#’s Parallel.ForEach method, which will be referred to for this type of operation.

Thread synchronization is not required when iterating, however keeping track of the thread state is important in some instances.

* 1. **Index-Based Methods**

Accessing or modifying an element at a specified index (such as get, set, or removeAt methods) is somewhat more complex than in a regular list.

To get the chunk where the position would be located, divide the index by the chunk size and cast it to an integer: 𝑐ℎ𝑢𝑛𝑘 = 𝑖𝑛t(𝑖𝑛𝑑𝑒𝑥 / 𝑐ℎ𝑢𝑛𝑘𝑆𝑖𝑧𝑒)

To get the position in the chunk where the index would be, use modulo on the index by the chunk size: 𝑐ℎ𝑢𝑛𝑘𝑃𝑜𝑠𝑖𝑡𝑖𝑜𝑛 = 𝑖𝑛𝑑𝑒𝑥 % 𝑐ℎ𝑢𝑛𝑘𝑆𝑖𝑧𝑒

Access via list*[chunk][chunkPosition]* (Where list is the main list inside the class)

*Index-Accessing Example:*

private int convertIndexToChunk(int index)

{

return index / chunkSize;

}

private int convertIndexToChunkPos(int index)

{

return index % chunkSize;

}

*Index Example:*

*The following example demonstrates accessing an element at index 8 in a chunk list containing numbers 0 – 10 with chunk size 5. (See Figure 2)*

**Accessing the chunk:** int(8 / 5) = 1

**Accessing the chunk position:** 8 % 5 = 3

***Fig. 2:***

|  |  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
| **Chunk 0**   |  |  |  |  |  | | --- | --- | --- | --- | --- | | [0] | [1] | [2] | [3] | [4] | | 0 | 1 | 2 | 3 | 4 | |
| **Chunk 1**   |  |  |  |  |  | | --- | --- | --- | --- | --- | | [0] | [1] | [2] | [3] | [4] | | 5 | 6 | 7 | ***8*** | 9 | |
| **Chunk 1**   |  |  |  |  |  | | --- | --- | --- | --- | --- | | [0] | [1] | [2] | [3] | [4] | | 10 |  |  |  |  | |

### 2.3.1 Index Issues

One issue with using indices in a chunk list, however, is the problem where items flow left within the chunk but do not migrate left from one to another if a chunk has an open slot. To implement so may hinder performance during removal.

However, a very simple solution would be to use recursion, such as within a try-catch statement using the index + 1.

*Example:*

*The following example demonstrates a solution to the problem by counting up the index until an open position is found:*

public T get(int index)

{

if (index >= size()) throw new ArgumentOutOfRangeException();

try

{

return myList[convertIndexToChunk(index)][convertIndexToChunkPos(index)];

}

catch (ArgumentOutOfRangeException)

{

return get(index + 1);

}

}

**2.3.2 Chunk Resizing**

Should our data set grow marginally larger, we may need to resize our list. To do so however, means we’ll need to rebalance our list, which is especially important if the chunk size we’re changing to is smaller than the current one.

We can make a temporary list containing all of our old items, change the chunk size, clear our old list, and then reflow our data back in.

While somewhat costly performance-wise, this is an operation that should not be necessary to occur often.

If the chunk size we want to adjust to is larger than the current one, however, we can simply leave the list as is and allow the elements to re-fill the chunks that are not yet at capacity.

*Example:*

public void setChunkSize(int newChunkSize)

{

if (newChunkSize > chunkSize)

{

chunkSize = newChunkSize;

}

else

{

var items = getList();

chunkSize = newChunkSize;

clear();

foreach (var item in items)

{

add(items);

}

}

}

public List<T> getList()

{

var items = new List<T>();

foreach (var currentList in myList)

{

foreach (T currentItem in currentList)

{

items.Add(currentItem);

}

}

return items;

}

* 1. **Element Operations**

A CHUNK LIST IS AN ARRAY-BASED LIST OF ELEMENTS IN WHICH

**2.4.1 Adding Elements**

Adding elements to a chunk list is simple; however, it does require that we check if each chunk is at capacity. Getting the size from the chunk should be Big-O (1), so this should not increase runtime marginally whatsoever.

An element will naturally fall into the first open spot, or the first chunk that is not at capacity. (See Figure 3)

If all chunks are at capacity, however, we need to add a new chunk to our list, then add the item it.

***Fig. 3:***

|  |
| --- |
| **10** |

|  |  |  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
| |  |  |  |  |  | | --- | --- | --- | --- | --- | | 0 | 1 | 2 | 3 | 4 | | |  |  |  |  |  | | --- | --- | --- | --- | --- | | 5 | 6 | 7 | 8 | 9 | |

|  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- |
| **Chunk 0**   |  |  |  |  |  | | --- | --- | --- | --- | --- | | 0 | 1 | 2 | 3 | 4 | |
| **Chunk 1**   |  |  |  |  |  | | --- | --- | --- | --- | --- | | 5 | 6 | 7 | 8 | 9 | |
| **Chunk 1**   |  |  |  |  |  | | --- | --- | --- | --- | --- | | **10** |  |  |  |  | |

*Example:*

public void add(T t)

{

foreach (List<T> currentList in myList)

{

if (currentList.Count != chunkSize)

{

currentList.Add(t);

return;

}

}

myList.Add(new List<T>());

myList[myList.Count – 1].Add(t);

}

**2.4.2 Removing Elements**

Removing elements is one of the fastest computational operations in a chunk list. This is where we can start using multithreading to our advantage.

To remove an element, we can use a parallel for loop to concurrently check each chunk for the item.

We can use a binary search to get the index that we’re looking for.

This is also where we need to be able to have access to the thread’s state when we’re looping through each chunk. If we only want to remove the first found instance of an element, we need to immediately break out of the parallel for loop.

To remove all instances of an element within the list, we can still use a parallel for loop, and just call a removeAll method on each chunk.

To clear the entire list, we can simply call clear on the main list (containing the chunks).

*Example:*

public void remove(T t)

{

Parallel.ForEach(myList, (currentList, state) =>

{

int indx = currentList.BinarySearch(t);

if (indx >= 0)

{

currentList.RemoveAt(indx);

state.Break();

}

});

}

public void removeAll(T t)

{

Parallel.ForEach(myList, (currentList) =>

{

for (int i = 0; i < currentList.Count; i++)

{

if (currentList[i].Equals(t))

{

currentList.RemoveAt(i);

i--;

}

}

});

}

**2.4.3 Searching**

Searching for an element is also where chunk lists shine. Once again we can use concurrency to get the shortest possible runtime, as now we can use a parallel for loop not only on the list itself, but on each chunk.

Essentially, we can check most items in the list at the exact same time, meaning our runtime will be marginally smaller than using a linear search at worst case, and in the best case, a binary search.

*Example:*

public bool contains(T t)

{

bool found = false;

Parallel.ForEach(myList, (currentList, state) =>

{

Parallel.ForEach(currentList, (currentItem) =>

{

if (currentItem,Equals(t))

{

found = true;

state.Break();

}

});

});

return found;

}

**2.4.4 Sorting**

Sorting our list is a fairly complex operation, similarly to searching. To properly sort our list, we do have to make a temporary list containing all elements of our chunk list. To do otherwise would only sort the chunks, which is not ideal as we do not know which order they will be inserted in.

Using our temporary list, we can clear our main list and simply reflow all of our items back in after sorting it.

*Sorting Example:*

*For this example, I simply used the sort method implemented within C#’s List class, which follows the following rules:*

* *If the partition size is fewer than 16 elements, it uses an insertion sort algorithm.*
* *If the number of partitions exceeds 2 \* log N, where N is the range of the input array, it uses a heapsort algorithm.*
* *Otherwise, it uses a quicksort algorithm.*

public void sort()

{

var items = getList();

items.Sort();

clear();

foreach (T item in items)

{

add(item);

}

}

**3. COMPLEXITY**

**3.1 Complexities – Basic Methods**

***Complexities are listed with the following variables:***

* C being the number of chunks currently in the list.
* N being the number of elements per chunk.
* P being the number of processors.
* I being the index input for the operation.

|  |  |  |
| --- | --- | --- |
| **OPERATION** | **AVERAGE CASE** | **WORST CASE** |
| Add | Θ(C) | N/A |
| Remove | Θ((log C \* log N) / P) | Θ(log C \* log N) |
| RemoveAll | Θ((log C \* N) / P) | Θ(log C \* N) |
| RemoveAt | Θ(1) | Θ(C \* N - I) |
| Set | Θ(1) | Θ(C \* N - I) |
| Get | Θ(1) | Θ(C \* N - I) |

**3.2 Complexities – Additional Methods**

|  |  |  |
| --- | --- | --- |
| **OPERATION** | **AVERAGE CASE** | **WORST CASE** |
| GetList | Θ(C ^ 2 \* N) | N/A |
| Contains (Search) | Θ((log C \* log N) / P) | Θ(C \* N) |
| Size (Count) | Θ(C) | N/A |
| SetChunkSize | Θ(1) | Θ(C ^ 2 \* N) |
| Sort | Θ(C \* N \* log N) | Θ(C \* N ^ 2) |

**4. USAGE**

**4.1 Modern Usage**

A CHUNK LIST IS AN ARRAY-BASED LIST OF ELEMENTS IN WHICH

**4.1.1 Experiments**

A CHUNK LIST IS AN ARRAY-BASED LIST OF ELEMENTS IN WHICH

**4.1.2 Unit Test & Benchmarks**

A CHUNK LIST IS AN ARRAY-BASED LIST OF ELEMENTS IN WHICH

**4.2 Related Work**

A CHUNK LIST IS AN ARRAY-BASED LIST OF ELEMENTS IN WHICH

**REFERENCES**

**APPENDIX**

*Method Headers & Complexities*

|  |  |  |
| --- | --- | --- |
| **OPERATION** | **AVERAGE CASE** | **WORST CASE** |
| **Constructor** *()* *:* *this(DEFAULT\_SIZE)* | *N/A* | *N/A* |
| **Constructor** *(int chunkSize)* | *N/A* | *N/A* |
| void **Add** *(T t)* | Θ(C) | *N/A* |
| void **Remove** *(T t)* | Θ((log C \* log N) / P) | Θ(log C \* log N) |
| void **RemoveAll** *(T t)* | Θ((log C \* N) / P) | Θ(log C \* N) |
| void **RemoveAt** *(int index)* | Θ(1) | Θ(C \* N - I) |
| int **ConvertIndexToChunk** *(int index)* | *N/A* | *N/A* |
| int **ConvertIndexToChunkPos** *(int index)* | *N/A* | *N/A* |
| void **Set** *(int index, T t)* | Θ(1) | Θ(C \* N - I) |
| T **Get** *(int index)* | Θ(1) | Θ(C \* N - I) |
| List<T> **GetList** *()* | Θ(C ^ 2 \* N) | N/A |
| bool **Contains** (Search) *(T t)* | Θ((log C \* log N) / P) | Θ(C \* N) |
| int **Size** (Count) *()* | Θ(C) | *N/A* |
| void **SetChunkSize** *(int size)* | Θ(1) | Θ(C ^ 2 \* N) |
| void **Sort** *()* | Θ(C \* N \* log N) | Θ(C \* N ^ 2) |

© Daniel Szelogowski 2017